

Inexact shift-and-invert Arnoldi's method and implicit restarts with preconditioning for eigencomputations

Melina Freitag

Department of Mathematical Sciences
University of Bath

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joint work with Howard Elman (Maryland) and Alastair Spence (Bath)

- 1 Introduction
- 2 Inexact Shift-invert Arnoldi method
- 3 Inexact Shift-invert Arnoldi method with implicit restarts
- 4 Conclusions

Outline

1 Introduction

2 Inexact Shift-invert Arnoldi method

3 Inexact Shift-invert Arnoldi method with implicit restarts

4 Conclusions

Problem and iterative methods

Find a small number of eigenvalues and eigenvectors of:

$$Ax = \lambda x, \quad \lambda \in \mathbb{C}, x \in \mathbb{C}^n$$

- A is large, sparse, **nonsymmetric**

Problem and iterative methods

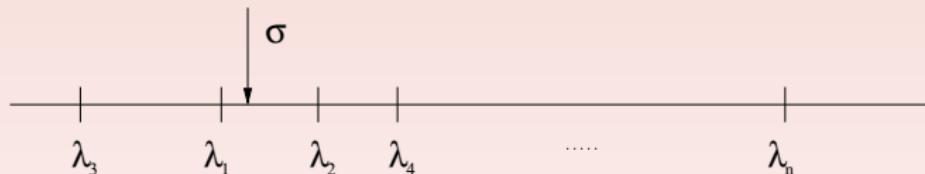
Find a small number of eigenvalues and eigenvectors of:

$$Ax = \lambda x, \quad \lambda \in \mathbb{C}, x \in \mathbb{C}^n$$

- A is large, sparse, **nonsymmetric**
- Iterative solves
 - Power method
 - Simultaneous iteration
 - **Arnoldi method**
 - Jacobi-Davidson method
- The first three of these involve repeated application of the matrix A to a vector
- Generally convergence to largest/outlying eigenvector

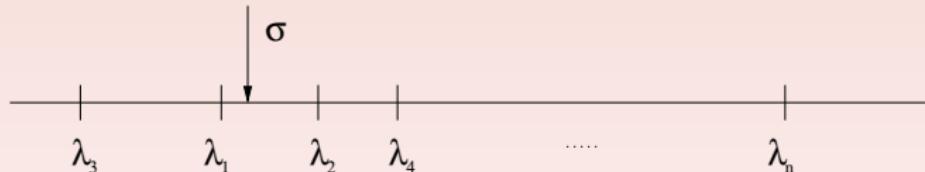
Shift-invert strategy

- Wish to find a few eigenvalues close to a shift σ



Shift-invert strategy

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- Problem becomes

$$(A - \sigma I)^{-1}x = \frac{1}{\lambda - \sigma}x$$

- each step of the iterative method involves repeated application of $(A - \sigma I)^{-1}$ to a vector

- Inner iterative solve:**

$$(A - \sigma I)y = x$$

using Krylov method for linear systems.

- leading to **inner-outer iterative method**.

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The algorithm

Arnoldi's method

- Arnoldi method constructs an orthogonal basis of k -dimensional Krylov subspace

$$\mathcal{K}_k(\mathcal{A}, q^{(1)}) = \text{span}\{q^{(1)}, \mathcal{A}q^{(1)}, \mathcal{A}^2q^{(1)}, \dots, \mathcal{A}^{k-1}q^{(1)}\},$$

$$\mathcal{A}Q_k = Q_k H_k + q_{k+1} h_{k+1,k} e_k^H = Q_{k+1} \begin{bmatrix} H_k \\ h_{k+1,k} e_k^H \end{bmatrix}$$

$$Q_k^H Q_k = I.$$

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- Eigenvalues of H_k are eigenvalue approximations of (outlying) eigenvalues of \mathcal{A}

$$\|r_k\| = \|\mathcal{A}x - \theta x\| = \|(\mathcal{A}Q_k - Q_k H_k)u\| = |h_{k+1,k}| |e_k^H u|,$$

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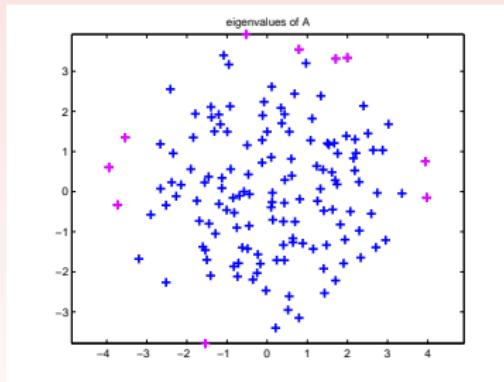
$$\|r_k\| = \|\mathcal{A}x - \theta x\| = \|(\mathcal{A}Q_k - Q_k H_k)u\| = |h_{k+1,k}| |e_k^H u|,$$

- at each step, application of \mathcal{A} to q_k : $\mathcal{A}q_k = \tilde{q}_{k+1}$

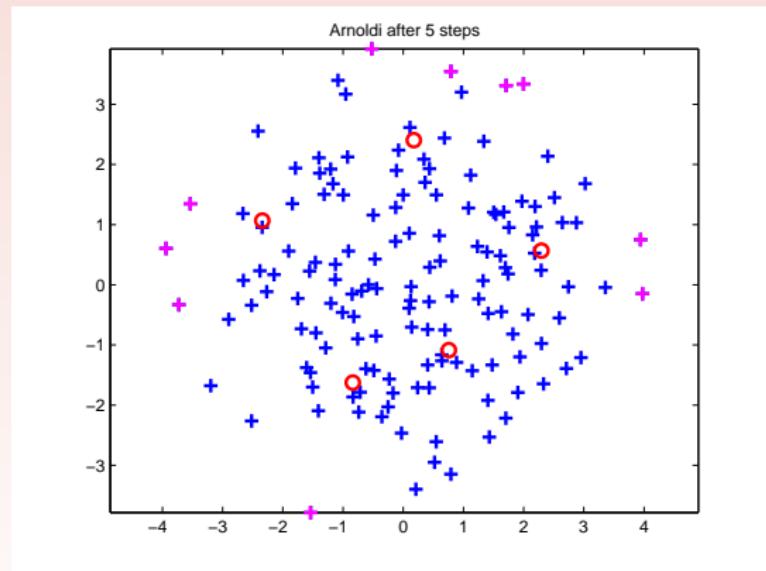
Example

random complex matrix of dimension $n = 144$ generated in MATLAB:

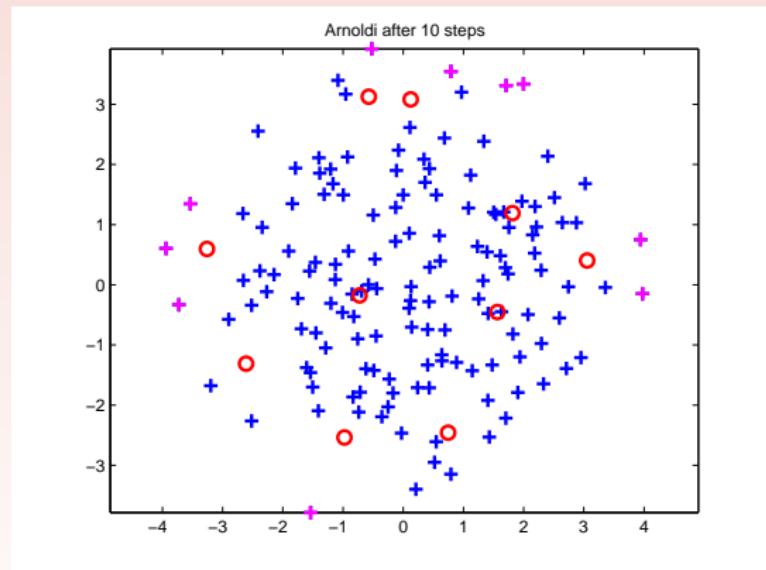
```
G=numgrid('N',14);B=delsq(G);A=sprandn(B)+i*sprandn(B)
```



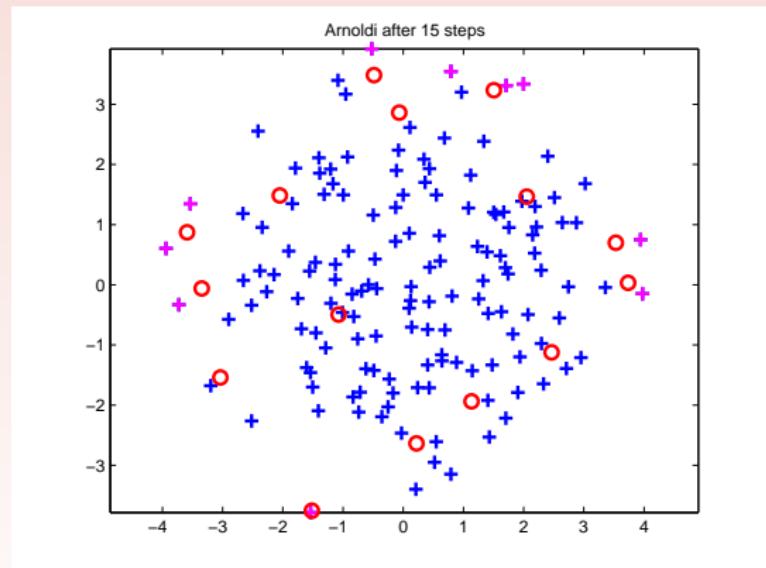
after 5 Arnoldi steps



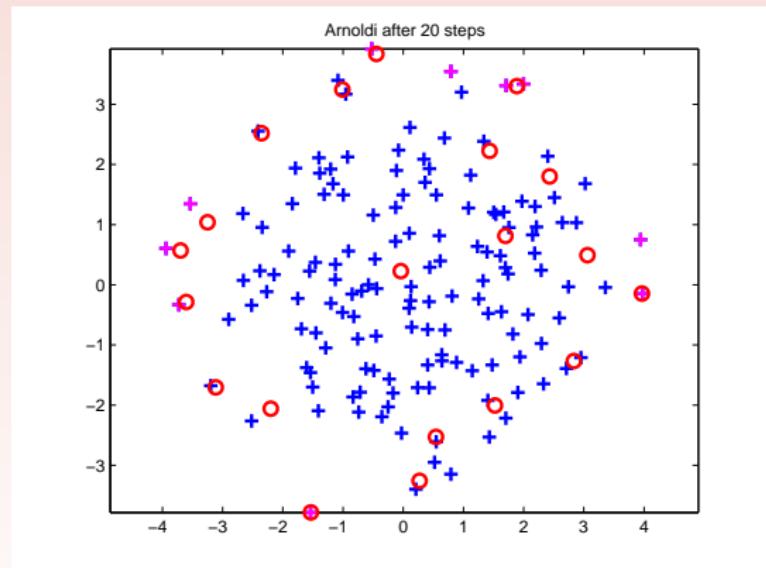
after 10 Arnoldi steps



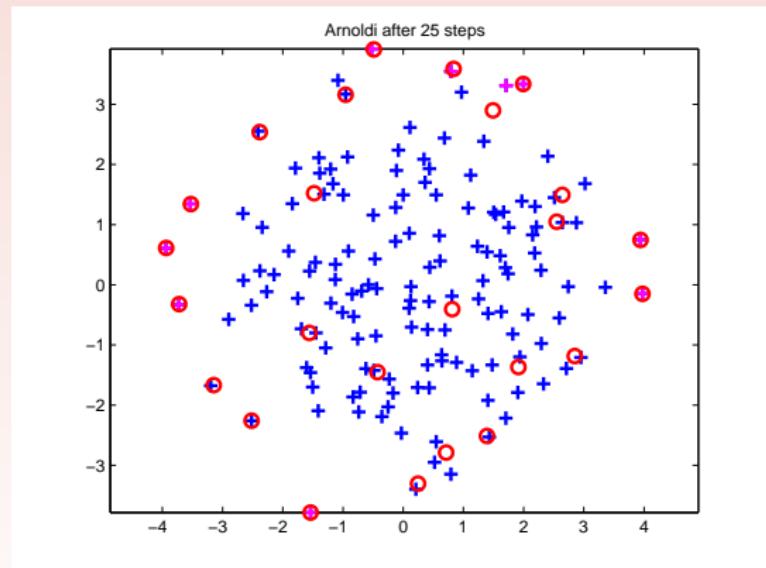
after 15 Arnoldi steps



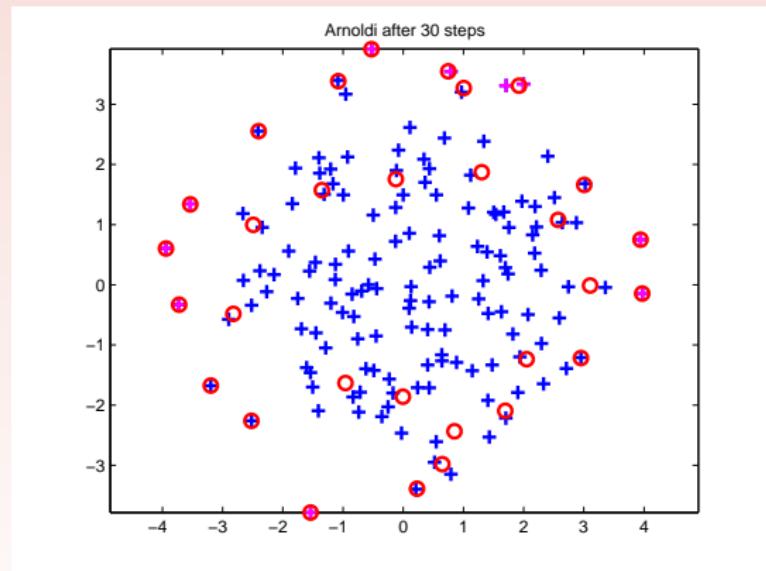
after 20 Arnoldi steps



after 25 Arnoldi steps



after 30 Arnoldi steps



The algorithm: take $\sigma = 0$

Shift-Invert Arnoldi's method $\mathcal{A} := A^{-1}$

- Arnoldi method constructs an orthogonal basis of k -dimensional Krylov subspace

$$\mathcal{K}_k(A^{-1}, q^{(1)}) = \text{span}\{q^{(1)}, A^{-1}q^{(1)}, (A^{-1})^2q^{(1)}, \dots, (A^{-1})^{k-1}q^{(1)}\},$$

$$A^{-1}Q_k = Q_k H_k + q_{k+1} h_{k+1,k} e_k^H = Q_{k+1} \begin{bmatrix} H_k \\ h_{k+1,k} e_k^H \end{bmatrix}$$

$$Q_k^H Q_k = I.$$

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$$\|r_k\| = \|A^{-1}x - \theta x\| = \|(A^{-1}Q_k - Q_k H_k)u\| = |h_{k+1,k}| |e_k^H u|,$$

- at each step, application of A^{-1} to q_k : $A^{-1}q_k = \tilde{q}_{k+1}$

Inexact solves

Inexact solves (Simoncini 2005), Bouras and Frayssé (2000)

- Wish to solve

$$\|q_k - A\tilde{q}_{k+1}\| = \|\tilde{d}_k\| \leq \tau_k$$

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- leads to **inexact Arnoldi relation**

$$A^{-1}Q_k = Q_{k+1} \begin{bmatrix} H_k \\ h_{k+1,k}e_k^H \end{bmatrix} + \textcolor{red}{D}_k = Q_{k+1} \begin{bmatrix} H_k \\ h_{k+1,k}e_k^H \end{bmatrix} + [d_1 | \dots | d_k]$$

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- u eigenvector of H_k :

$$\|r_k\| = \|(A^{-1}Q_k - Q_k H_k)u\| = |h_{k+1,k}| |e_k^H u| + \mathcal{D}_k u,$$

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$$\|r_k\| = \|(A^{-1}Q_k - Q_k H_k)u\| = |h_{k+1,k}| |e_k^H u| + \mathbf{D}_k u,$$

- Linear combination of the columns of D_k

$$\mathbf{D}_k u = \sum_{l=1}^k d_l u_l, \quad \text{if } u_l \text{ small, then } d_l \text{ allowed to be large!}$$

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$$\|d_l u_l\| \leq \frac{1}{k} \varepsilon \Rightarrow \|\mathbf{D}_k u\| < \varepsilon$$

and

$$|u_l| \leq C(l, k) \|r_{l-1}\| \quad \star$$

leads to

$$\|q_k - A\tilde{q}_{k+1}\| = \|\tilde{d}_k\|$$

$$\|\tilde{d}_k\| = C \frac{1}{\|r_{k-1}\|} \quad \diamond$$

The inner iteration for $AP^{-1}\tilde{q}_{k+1} = q_k$

Preconditioning

- Introduce preconditioner P and solve

$$AP^{-1}\tilde{q}_{k+1} = q_k, \quad P^{-1}\tilde{q}_{k+1} = q_{k+1}$$

using GMRES

The inner iteration for $(A - \sigma I)y = x$

GMRES

The k th iterate of GMRES applied to $Bz = b$ is the solution to

$$\min_{z_k \in z_0 + \mathcal{K}_k(B, r_0)} \|b - Bz_k\|,$$

where $\mathcal{K}_k(B, r_0) = \text{span}\{r_0, Br_0, B^2r_0, \dots, B^{k-1}r_0\}$. Set

$$z_k = z_0 + \sum_{j=0}^{k-1} \gamma_j B^j r_0$$

and

$$\|r_k\| = \|(I - \sum_{j=0}^{k-1} \gamma_j B^{j+1})r_0\| = \|p(B)r_0\|.$$

The inner iteration for $AP^{-1}\tilde{q}_{k+1} = q_k$

Preconditioning

■ GMRES convergence bound

$$\|d_l\| = \kappa \min_{p \in \Pi_l} \max_{i=1, \dots, n} |p(\mu_i)| \|d_0\|$$

depending on

The inner iteration for $AP^{-1}\tilde{q}_{k+1} = q_k$

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depending on

- the eigenvalue clustering of AP^{-1}
- the condition number
- the right hand side (initial guess)

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depending on

- the eigenvalue clustering of AP^{-1}
- the condition number
- the right hand side (initial guess)
- using a **tuned** preconditioner for Arnoldi's method

$$\mathbb{P}_k Q_k = A Q_k; \quad \text{given by} \quad \mathbb{P}_k = P + (A - P) Q_k Q_k^H$$

The inner iteration for $A\tilde{q} = q$

Theorem (Properties of the tuned preconditioner)

Let P with $P = A + E$ be a preconditioner for A and assume k steps of Arnoldi's method have been carried out; then k eigenvalues of $A\mathbb{P}_k^{-1}$ are equal to one:

$$[A\mathbb{P}_k^{-1}]AQ_k = AQ_k$$

and $n - k$ eigenvalues are close to the corresponding eigenvalues of AP^{-1} . They are eigenvalues of $L \in \mathbb{C}^{n-k \times n-k}$ with

$$\|L - I\| \leq C\|E\|.$$

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Implementation

- Sherman-Morrison-Woodbury.
- Only minor extra costs (one back substitution per outer iteration)

Numerical Example

`sherman5.mtx` nonsymmetric matrix from the Matrix Market library
(3312×3312).

- smallest eigenvalue: $\lambda_1 \approx 4.69 \times 10^{-2}$,
- Preconditioned GMRES as inner solver (both fixed tolerance and relaxation strategy),
- standard and tuned preconditioner (incomplete LU).

No tuning and standard preconditioner

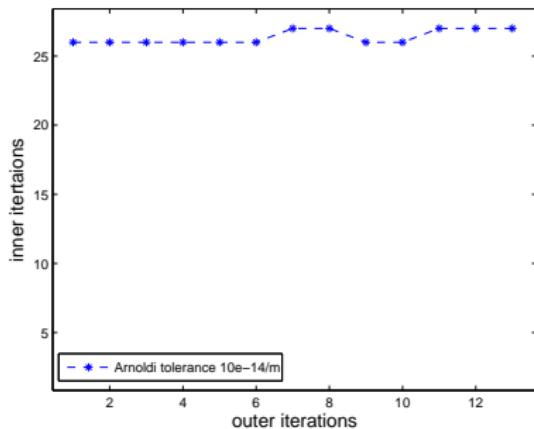


Figure: Inner iterations vs outer iterations

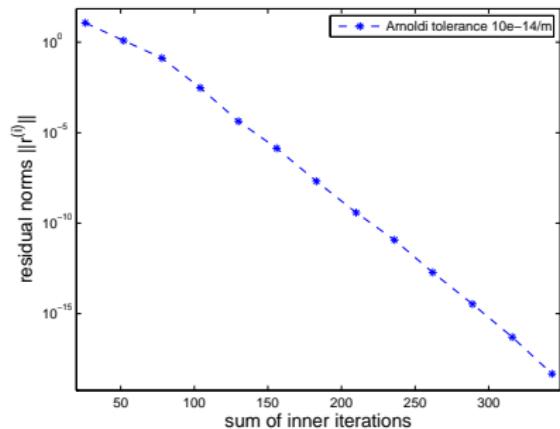


Figure: Eigenvalue residual norms vs total number of inner iterations

Tuning the preconditioner

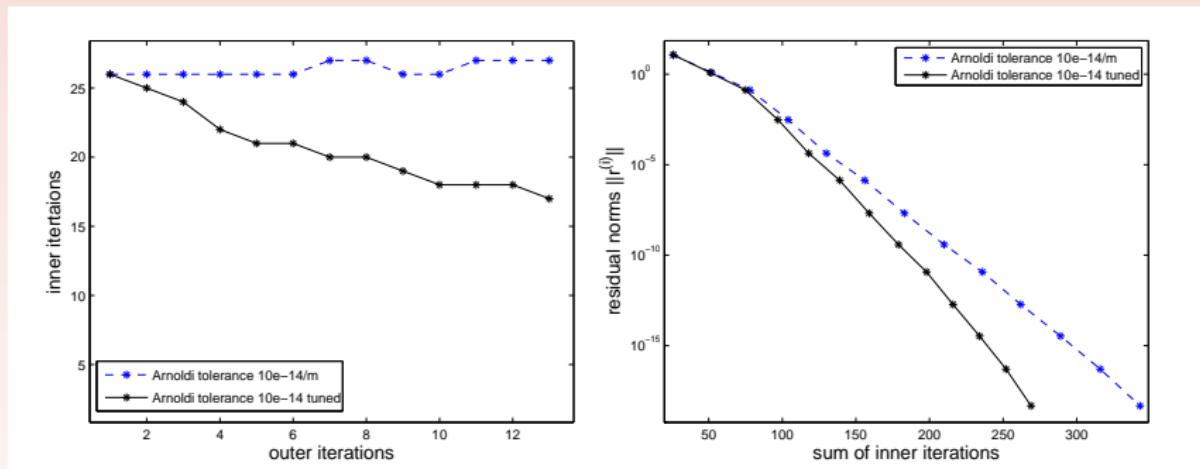


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Figure: Eigenvalue residual norms vs total number of inner iterations

Relaxation

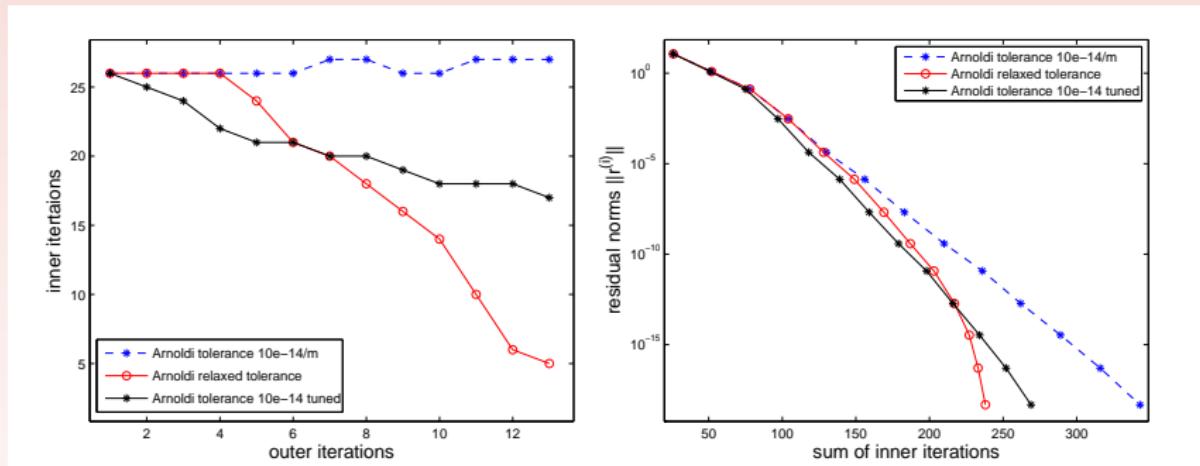


Figure: Inner iterations vs outer iterations

Figure: Eigenvalue residual norms vs total number of inner iterations

Tuning and relaxation strategy

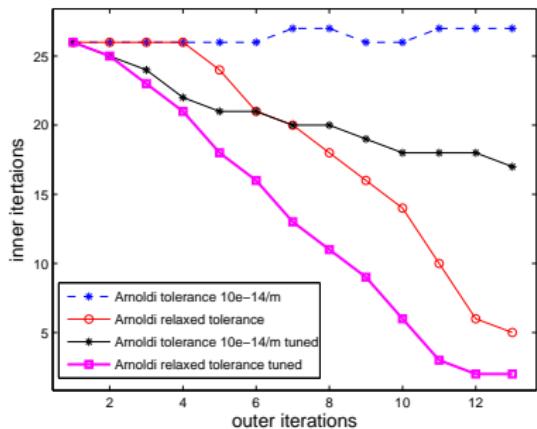


Figure: Inner iterations vs outer iterations

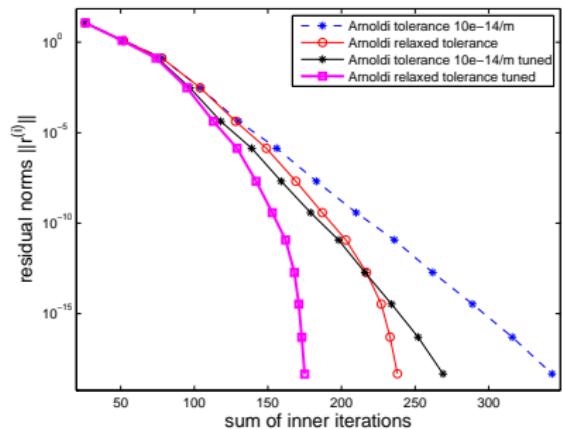


Figure: Eigenvalue residual norms vs total number of inner iterations

Ritz values of exact and inexact Arnoldi

Exact eigenvalues	Ritz values (exact Arnoldi)	Ritz values (inexact Arnoldi, tuning)
+4.69249563e-02	+ <u>4.69249563e-02</u>	+ <u>4.69249563e-02</u>
+1.25445378e-01	+ <u>1.25445378e-01</u>	+ <u>1.25445378e-01</u>
+4.02658363e-01	+ <u>4.02658347e-01</u>	+ <u>4.02658244e-01</u>
+5.79574381e-01	+ <u>5.79625498e-01</u>	+ <u>5.79817301e-01</u>
+6.18836405e-01	+ <u>6.18798666e-01</u>	+ <u>6.18650849e-01</u>

Table: Ritz values of exact Arnoldi's method and inexact Arnoldi's method with the tuning strategy compared to exact eigenvalues closest to zero after 14 shift-invert Arnoldi steps.

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Exact shifts

- take an $k + p$ step Arnoldi factorisation

$$\mathcal{A}Q_{k+p} = Q_{k+p}H_{k+p} + q_{k+p+1}h_{k+p+1,k+p}e_{k+p}^H$$

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- Compute $\Lambda(H_{k+p})$ and select p shifts for an implicit QR iteration

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- Compute $\Lambda(H_{k+p})$ and select p shifts for an implicit QR iteration
- implicit restart with new starting vector $\hat{q}^{(1)} = \frac{p(\mathcal{A})q^{(1)}}{\|p(\mathcal{A})q^{(1)}\|}$

Implicitly restarted Arnoldi (Sorensen (1992))

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Aim of IRA

$$\mathcal{A}Q_k = Q_kH_k + q_{k+1} \underbrace{h_{k+1,k}}_{\rightarrow 0} e_k^H$$

Relaxation strategy for IRA

Theorem

For any given $\varepsilon \in \mathbb{R}$ with $\varepsilon > 0$ assume that

$$\|d_l\| \leq \begin{cases} \varepsilon \frac{C}{\|R_k\|} & \text{if } l > k, \\ \varepsilon & \text{otherwise.} \end{cases} \quad \diamond$$

Then

$$\|\mathcal{A}Q_mU - Q_mU\Theta - R_m\| \leq \varepsilon.$$

- Very technical
- Relaxation strategy also works for IRA!

Tuning

Tuning for implicitly restarted Arnoldi's method

- Introduce preconditioner P and solve

$$A\mathbb{P}_k^{-1}\tilde{q}_{k+1} = q_k, \quad \mathbb{P}_k^{-1}\tilde{q}_{k+1} = q_{k+1}$$

using GMRES and a **tuned** preconditioner

$$\mathbb{P}_k Q_k = A Q_k; \quad \text{given by} \quad \mathbb{P}_k = P + (A - P) Q_k Q_k^H$$

Tuning

Why does tuning help?

- Assume we have found an approximate invariant subspace, that is

$$A^{-1}Q_k = Q_k H_k + \underbrace{q_{k+1} h_{k+1,k} e_k^H}_{\approx 0}$$

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- let A^{-1} have the upper Hessenberg form

$$[Q_k \ Q_k^\perp]^H \underbrace{A^{-1}}_{\text{Hessenberg}} [Q_k \ Q_k^\perp] = \begin{bmatrix} H_k & T_{12} \\ h_{k+1,k} e_k e_k^H & T_{22} \end{bmatrix},$$

where $[Q_k \ Q_k^\perp]$ is unitary and $H_k \in \mathbb{C}^{k,k}$ and $T_{22} \in \mathbb{C}^{n-k, n-k}$ are upper Hessenberg.

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If $h_{k+1,k} \neq 0$ then

$$[Q_k \ Q_k^\perp]^H \textcolor{brown}{A} \mathbb{P}_k^{-1} [Q_k \ Q_k^\perp] = \begin{bmatrix} I + \star & Q_k^H A \mathbb{P}_k^{-1} Q_k^\perp \\ \star & T_{22}^{-1} (Q_k^\perp)^H P Q_k^\perp + \star \end{bmatrix}$$

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where $[Q_k \ Q_k^\perp]$ is unitary and $H_k \in \mathbb{C}^{k,k}$ and $T_{22} \in \mathbb{C}^{n-k, n-k}$ are upper Hessenberg.

If $h_{k+1,k} = 0$ then

$$[Q_k \ Q_k^\perp]^H \underbrace{A \mathbb{P}_k^{-1}}_{\text{A} \mathbb{P}_k^{-1}} [Q_k \ Q_k^\perp] = \begin{bmatrix} I & Q_k^H A \mathbb{P}_k^{-1} Q_k^\perp \\ 0 & T_{22}^{-1} (Q_k^\perp)^H P Q_k^\perp \end{bmatrix}$$

Tuning

Another advantage of tuning

- System to be solved at each step of Arnoldi's method is

$$A\mathbb{P}_k^{-1}\tilde{q}_{k+1} = \textcolor{red}{q_k}, \quad \mathbb{P}_k^{-1}\tilde{q}_{k+1} = \tilde{q}_k$$

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- the right hand side of the system matrix is an eigenvector of the system!
- Krylov methods converge in one iteration

Tuning

Another advantage of tuning

- In practice:

$$A^{-1}Q_k = Q_k H_k + q_{k+1} h_{k+1,k} e_k^H$$

and

$$\|A\mathbb{P}_k^{-1}q_k - q_k\| = \mathcal{O}(|h_{k+1,k}|)$$

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- number of iterations decreases as the outer iteration proceeds

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$$A^{-1}Q_k = Q_k H_k + q_{k+1} h_{k+1,k} e_k^H$$

and

$$\|A\mathbb{P}_k^{-1}q_k - q_k\| = \mathcal{O}(|h_{k+1,k}|)$$

- number of iterations decreases as the outer iteration proceeds
- Rigorous analysis quite technical.

Numerical Example

`sherman5.mtx` nonsymmetric matrix from the Matrix Market library
(3312×3312).

- $k = 8$ eigenvalues closest to zero
- IRA with exact shifts $p = 4$
- Preconditioned GMRES as inner solver (fixed tolerance and relaxation strategy),
- standard and tuned preconditioner (incomplete LU).

No tuning and standard preconditioner

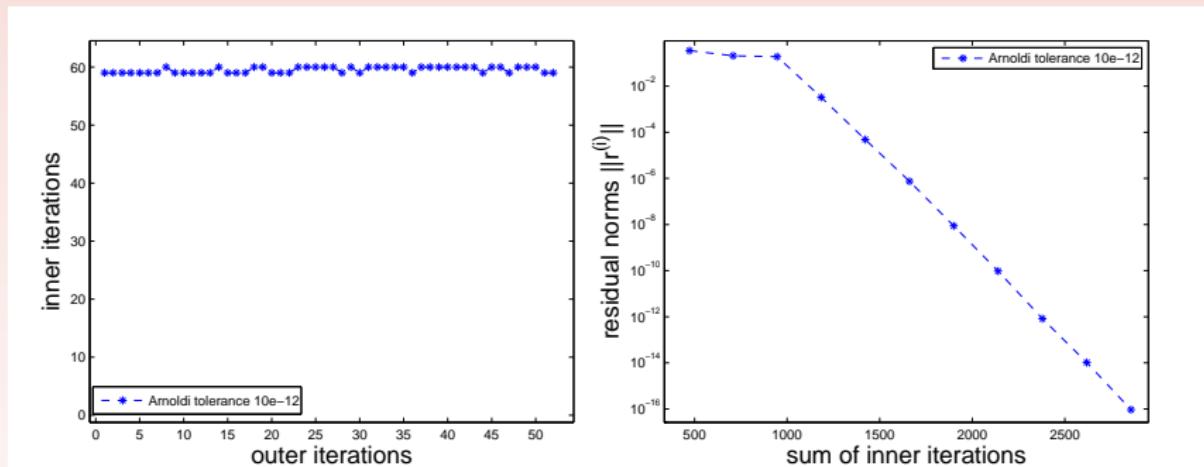


Figure: Inner iterations vs outer iterations

Figure: Eigenvalue residual norms vs total number of inner iterations

Tuning

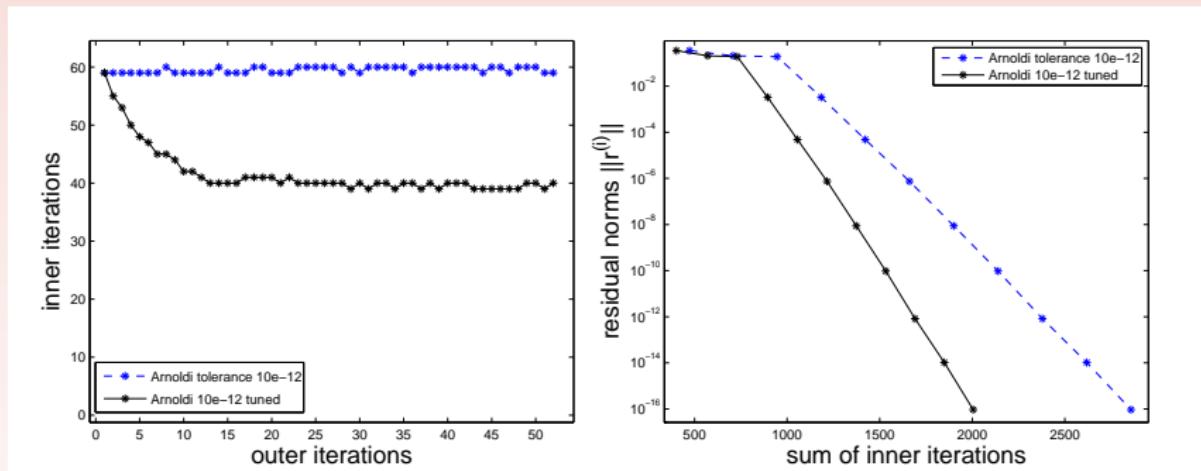


Figure: Inner iterations vs outer iterations

Figure: Eigenvalue residual norms vs total number of inner iterations

Relaxation

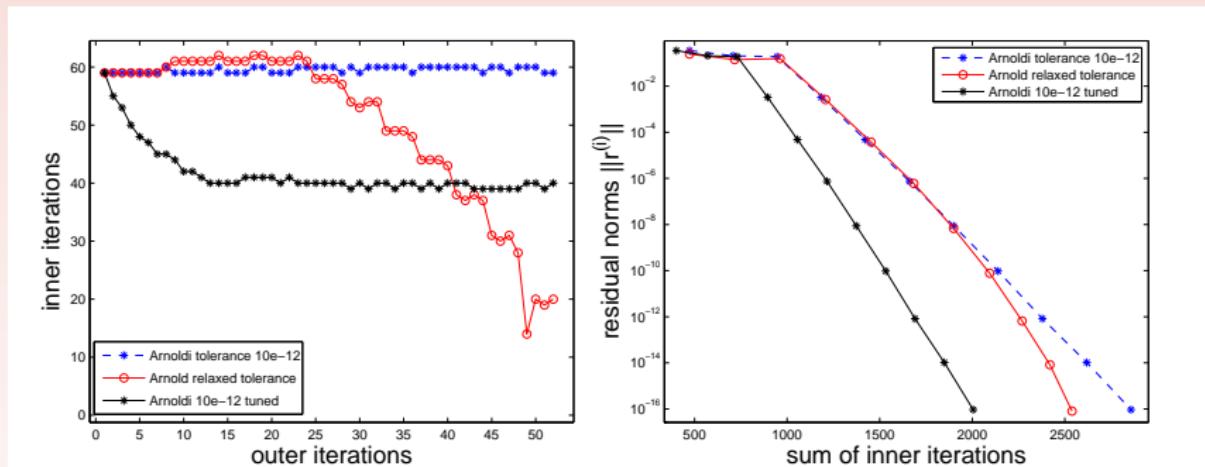


Figure: Inner iterations vs outer iterations

Figure: Eigenvalue residual norms vs total number of inner iterations

Tuning and relaxation strategy

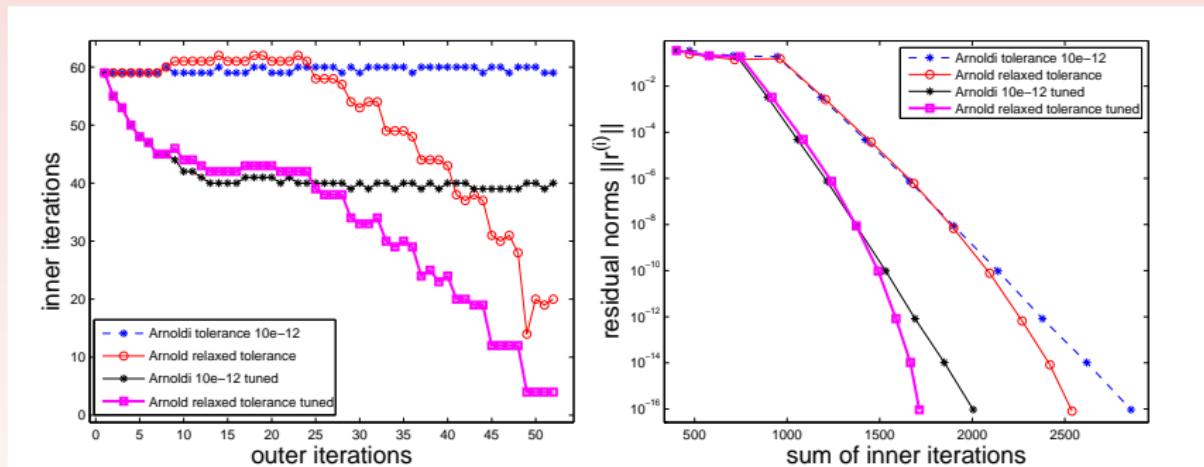


Figure: Inner iterations vs outer iterations

Figure: Eigenvalue residual norms vs total number of inner iterations

Numerical Example

`qc2534.mtx` matrix from the Matrix Market library.

- $k = 6$ eigenvalues closest to zero
- IRA with exact shifts $p = 4$
- Preconditioned GMRES as inner solver (fixed tolerance and relaxation strategy),
- standard and tuned preconditioner (incomplete LU).

Tuning and relaxation strategy

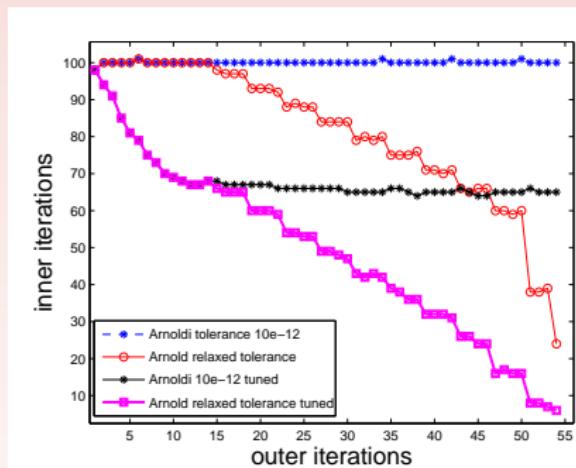


Figure: Inner iterations vs outer iterations

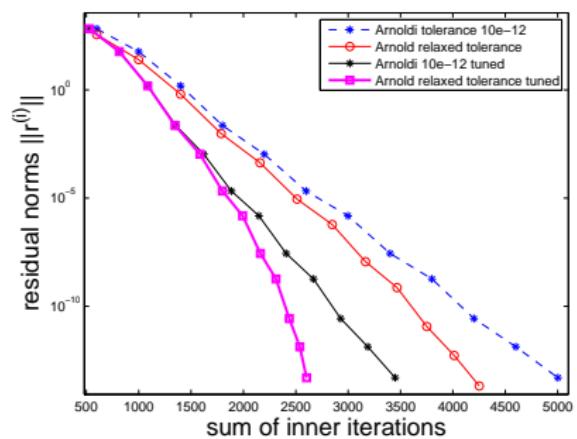


Figure: Eigenvalue residual norms vs total number of inner iterations

Outline

1 Introduction

2 Inexact Shift-invert Arnoldi method

3 Inexact Shift-invert Arnoldi method with implicit restarts

4 Conclusions

Conclusions

- For eigencomputations it is advantageous to consider small rank changes to the standard preconditioners (works for any preconditioner)
- Extension of the relaxation strategy to IRA
- Best results are obtained when relaxation and tuning are combined
- Current work: Link to Jacobi-Davidson



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